

# Quantum Annealing for Prime Factorization

Shuxian Jiang, Sabre Kais

Department of Chemistry and Computer Science, Purdue  
University

Keith A. Britt, Alexander J. McCaskey, Travis S. Humble  
Quantum Computing Institute, Oak Ridge National  
Laboratory

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# Integer Factorization Problem

## Classical Factorization

### - Methods:

1. Trial Division( $n^{\frac{1}{2}}$ ): try 2,3,5,..., $n^{\frac{1}{2}}$  to divide n
2. Shanks, Pollard, Strassen (1970s)( $n^{\frac{1}{4}}$ ): complicated math
3. CFRAC(1970), QS(1980), NFS(1990)  $O(\exp\sqrt{(\ln(n))(\ln(\ln(n)))})$ : nonrigorous
4. Elliptic Curve Method(1985)  $O(\exp\sqrt{(\ln(n))(\ln(\ln(n)))})$ : probabilistic, nonrigorous

### - Limitations:

**EXP** time if no randomness, no unproved hypotheses  
Quadratic Sieve Method(subexponential, not polynomial)

# Integer Factorization Problem

## Quantum Factorization

- Gate Model

Shor's algo: order-finding

Exponential speedup: polynomial time

Hard to physically implement

Largest Number: 21

- Adiabatic Computation Model

Convert to **optimization** problem

1. NMR

$$f = (N - pq)^2$$

2. Ising Machine: D-wave, ...

Multiplication Table

3. Nitrogen-Vacancy (NV) center in diamond

Multiplication Table

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## Quantum Factorization

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# Quantum Adiabatic Computation

## Method

- 1 Define Hamiltonian:

$$\mathcal{H}(t) = (1 - t/T)\mathcal{H}_0 + (t/T)\mathcal{H}_P$$

- 2 Set system to ground state of **initial**  $\mathcal{H}_0$  (easy)
- 3\* Define **final**  $\mathcal{H}_P$  s.t. its ground state encodes problem's solution
- 4 System slowly evolves to final state (eigenvector w.r.t. smallest eigenvalue), measure it to get soln

How to encode  $H_p$ 

$$f_p(s_1, s_2, \dots, s_n) = \sum_{k=1}^K \sum_{j_1, \dots, j_k=1}^n J_{j_1 \dots j_k} s_{j_1} \dots s_{j_k}, \quad s_i = \pm 1$$

is the energy func.(eigenvalue func.) of Hamiltonian

$$\mathcal{H}_p(\sigma_z^{(1)}, \sigma_z^{(2)}, \dots, \sigma_z^{(n)}) = \sum_{k=1}^K \sum_{j_1, \dots, j_k=1}^n J_{j_1 \dots j_k} \sigma_z^{(j_1)} \dots \sigma_z^{(j_k)}$$

Here  $\sigma_z^{(i)} = \overbrace{I \otimes I \otimes \dots \otimes I}^{i-1} \otimes \sigma_z \otimes \overbrace{I \otimes \dots \otimes I}^{n-i}$  with eigenvector  $|x_1 x_2 \dots x_n\rangle$ ,  $x_i = \{0, 1\}$ ,  $s_i = (-1)^{x_i}$

# Ising model

## Definition

$$\mathcal{H}_p(\sigma_z^{(1)}, \sigma_z^{(2)}, \dots, \sigma_z^{(n)}) = \sum_{i=1}^n h_i \sigma_z^{(i)} + \sum_{i,j=1}^n J_{ij} \sigma_z^{(i)} \sigma_z^{(j)}$$

$h_i$ : external magnetic

$J_{ij}$ : interaction between two adjacent sites  $i, j$

Ising Model  $\leftrightarrow$  Quadratic Function

# Our Method for Factorization<sup>1</sup>

- Define Cost Function

- Direct Method:  $f = (N - pq)^2$
- Modified Multiplication Table Method

- Order Reducing Method

- $x, y, z \in \{0, 1\}$

$$xy = z \text{ iff } xy - 2xz - 2yz + 3z = 0,$$

$$xy \neq z \text{ iff } xy - 2xz - 2yz + 3z > 0$$

$$\Rightarrow x_1, x_2, x_3 \in \{0, 1\}$$

$$\min(x_1 x_2 x_3) = \min_{x_4=x_1 x_2} (x_4 x_3 + 2(x_1 x_2 - 2x_1 x_4 - 2x_2 x_4 + 3x_4))$$

$$\min(-x_1 x_2 x_3) = \min_{x_4=x_1 x_2} (-x_4 x_3 + 2(x_1 x_2 - 2x_1 x_4 - 2x_2 x_4 + 3x_4))$$

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<sup>1</sup>Shuxian Jiang et al. "Quantum Annealing for Prime Factorization". In: *arXiv preprint arXiv:1804.02733* (2018).



## Cost Function: Direct Method

For  $N = pq$

$$p = (p_{l-1}p_{l-2}\dots p_1 1)_2, q = (q_{l'-1}q_{l'-2}\dots q_1 1)_2$$

$$\begin{aligned} \text{Define } f &= (N - pq)^2 \\ &= \left[ N - \left( \sum_{i=1}^{l-1} 2^i p_i + 1 \right) \left( \sum_{j=1}^{l'-1} 2^j q_j + 1 \right) \right]^2 \\ &= N^2 + \left( \sum_{i=1}^{l-1} 2^i p_i + 1 \right)^2 \left( \sum_{j=1}^{l'-1} 2^j q_j + 1 \right)^2 \\ &\quad - 2N \left( \sum_{i=1}^{l-1} 2^i p_i + 1 \right) \left( \sum_{j=1}^{l'-1} 2^j q_j + 1 \right) \end{aligned}$$

Reduce Order:

$$\min(x_1 x_2 x_3) = \min_{x_4=x_1 x_2} (x_4 x_3 + 2(x_1 x_2 - 2x_1 x_4 - 2x_2 x_4 + 3x_4))$$

## Cost Function: Direct Method

For  $N = 15 = 5 \times 3$

$$p = (x_1 1)_2 = x_1 * 2 + 1, q = (x_2 x_3 1)_2 = x_2 * 2^2 + x_3 * 2 + 1$$

$$\begin{aligned} & f(x_1, x_2, x_3) \\ &= (N - pq)^2 \\ &= [15 - (x_1 * 2 + 1)(x_3 * 2^2 + x_2 * 2 + 1)]^2 \\ &= 15^2 + (2x_1 + 1)^2(4x_3 + 2x_2 + 1)^2 - 30(2x_1 + 1)(4x_3 + 2x_2 + 1) \\ &= 128x_1x_2x_3 - 56x_1x_2 - 48x_1x_3 + 16x_2x_3 - 52x_1 - 52x_2 - 96x_3 \\ &\quad + 196 \end{aligned}$$

Reduce Order:

$$\begin{aligned} & f'(x_1, x_2, x_3, x_4) \\ &= 128(x_4x_3 + 2(x_1x_2 - 2x_1x_4 - 2x_2x_4 + 3x_4)) - 56x_1x_2 \\ &\quad - 48x_1x_3 + 16x_2x_3 - 52x_1 - 52x_2 - 96x_3 + 196 \end{aligned}$$

# Cost Function: Modified Multiplication Table Method

Table: Multiplication table for  $11 \times 13 = 143$  in binary.

	$2^7$	$2^6$	$2^5$	$2^4$	$2^3$	$2^2$	$2^1$	$2^0$
p					1	$p_2$	$p_1$	1
q					1	$q_2$	$q_1$	1
					1	$p_2$	$p_1$	1
				$q_1$	$p_2 q_1$	$p_1 q_1$	$q_1$	
		$q_2$	$p_2 q_2$	$p_1 q_2$	$q_2$			
		1	$p_2$	$p_1$	1			
carries		$c_4$	$c_3$	$c_2$	$c_1$			
$p \times q = 143$	1	0	0	0	1	1	1	1

$$2(p_2 + p_1 q_1 + q_2) + (p_1 + q_1) = 8c_2 + 4c_1 + 3$$

...

## Cost Function: Modified Multiplication Table Methods

From

$$2(p_2 + p_1 q_1 + q_2) + (p_1 + q_1) = 8c_2 + 4c_1 + 3$$

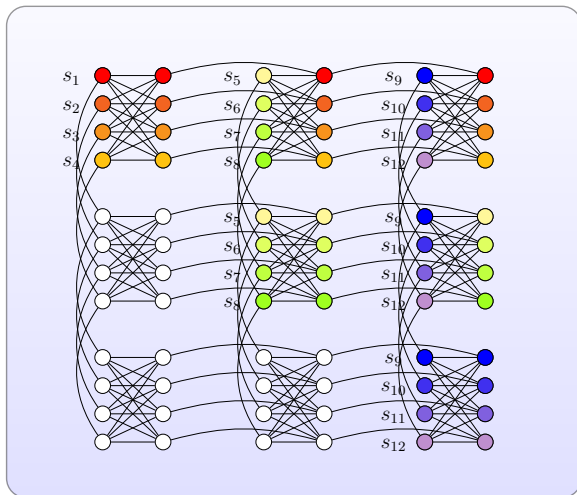
We got

$$f_1 = (2p_2 + 2p_1 q_1 + 2q_2 - 8c_2 - 4c_1 + p_1 + q_1 - 3)^2$$

Reduce order

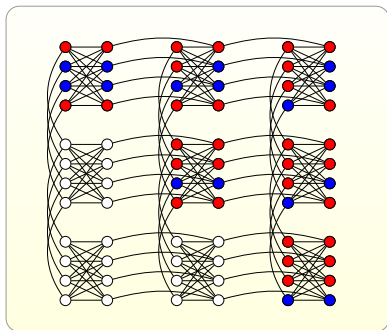
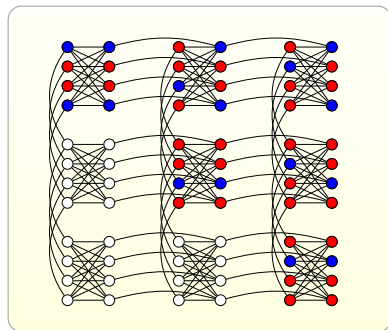
$$\min(\pm x_1 x_2 x_3) = \min_{x_4 = x_1 x_2} (\pm x_4 x_3 + 2(x_1 x_2 - 2x_1 x_4 - 2x_2 x_4 + 3x_4))$$

## Embed to D-wave



**Figure:** Embed problem graph of factoring 143 using method one to Chimera graph

## Embed to D-wave

(a) Result  $13 \times 11$ (b) Result  $11 \times 13$ 

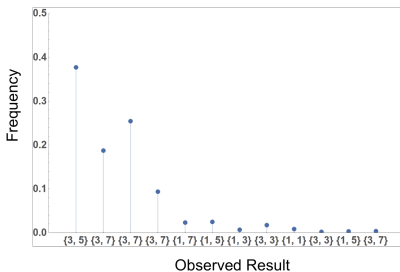
**Figure:** Final ground state of factoring 143. Red nodes: +1, Blue nodes: -1. (a)  $s_1 = 1, s_2 = -1, s_3 = -1, s_4 = 1$  ( $p = 1101, q = 1011$ ). (b)  $s_1 = -1, s_2 = 1, s_3 = 1, s_4 = -1$  ( $p = 1011, q = 1101$ ).

# Cost Function: Modified Multiplication Table Method

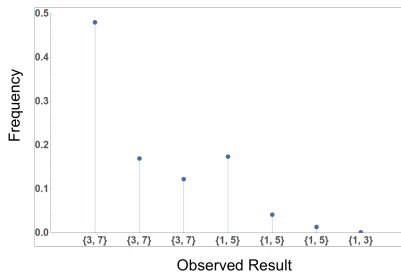
Table: Multiplication table for  $659 \times 571 = 376289$  in binary.

$2^{18}$	$2^{17}$	$2^{16}$	$2^{15}$	$2^{14}$	$2^{13}$	$2^{12}$	$2^{11}$	$2^{10}$	$2^9$	$2^8$	$2^7$	$2^6$	$2^5$	$2^4$	$2^3$	$2^2$	$2^1$	$2^0$
									1	$p_8$	$p_7$	$p_6$	$p_5$	$p_4$	$p_3$	$p_2$	$p_1$	1
									1	$q_8$	$q_7$	$q_6$	$q_5$	$q_4$	$q_3$	$q_2$	$q_1$	1
									1	$p_8$	$p_7$	$p_6$	$p_5$	$p_4$	$p_3$	$p_2$	$p_1$	1
								$q_1$	$p_8 q_1$	$p_7 q_1$	$p_6 q_1$	$p_5 q_1$	$p_4 q_1$	$p_3 q_1$	$p_2 q_1$	$p_1 q_1$	$q_1$	
							$q_2$	$p_8 q_2$	$p_7 q_2$	$p_6 q_2$	$p_5 q_2$	$p_4 q_2$	$p_3 q_2$	$p_2 q_2$	$p_1 q_2$	$q_2$		
					$q_3$	$p_8 q_3$	$p_7 q_3$	$p_6 q_3$	$p_5 q_3$	$p_4 q_3$	$p_3 q_3$	$p_2 q_3$	$p_1 q_3$	$q_3$				
				$q_4$	$p_8 q_4$	$p_7 q_4$	$p_6 q_4$	$p_5 q_4$	$p_4 q_4$	$p_3 q_4$	$p_2 q_4$	$p_1 q_4$	$q_4$					
			$q_5$	$p_8 q_5$	$p_7 q_5$	$p_6 q_5$	$p_5 q_5$	$p_4 q_5$	$p_3 q_5$	$p_2 q_5$	$p_1 q_5$	$q_5$						
		$q_6$	$p_8 q_6$	$p_7 q_6$	$p_6 q_6$	$p_5 q_6$	$p_4 q_6$	$p_3 q_6$	$p_2 q_6$	$p_1 q_6$	$q_6$							
	$q_7$	$p_8 q_7$	$p_7 q_7$	$p_6 q_7$	$p_5 q_7$	$p_4 q_7$	$p_3 q_7$	$p_2 q_7$	$p_1 q_7$	$q_7$								
	$q_8$	$p_8 q_8$	$p_7 q_8$	$p_6 q_8$	$p_5 q_8$	$p_4 q_8$	$p_3 q_8$	$p_2 q_8$	$p_1 q_8$	$q_8$								
1	$p_8$	$p_7$	$p_6$	$p_5$	$p_4$	$p_3$	$p_2$	$p_1$	1									
$c_{14}$	$c_{13}$			$c_9$	$c_8$	$c_7$	$c_6$	$c_5$	$c_4$	$c_3$			$c_2$	$c_1$				
1	0	1	1	0	1	1	1	1	0	1	1	1	1	0	0	0	0	1

# Experimental Results



(a)  $N = 15$ ,  $T = 200ms$



(b)  $N = 21$ ,  $T = 2000ms$

**Figure:** D-wave Results for Method One: rates of getting different solutions.



# Experimental Results

Table: D-wave Results for Method Two

N	factors	logic#qb	phy#qb	embed	sol
35	$7 \times 5$	5	8	✓	✓
143	$13 \times 11$	12	48	✓	✓
391	$23 \times 17$	20	113	✓	✓
1517	$41 \times 37$	30	270	✓	✓
8137	$103 \times 79$	43	535	✓	✓
56153	$241 \times 233$	58	1085	✓	—
249919	$509 \times 491$	77	1803	✓	—
376289	$659 \times 571$	94	—	—	—